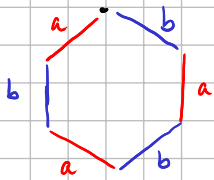


COMBINATORIAL TOPOLOGY AND GROUP THEORY

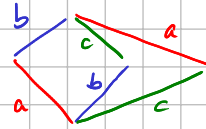
Lecture IV
03/12/2021

$(\mathcal{O}_3, \text{weak order})$



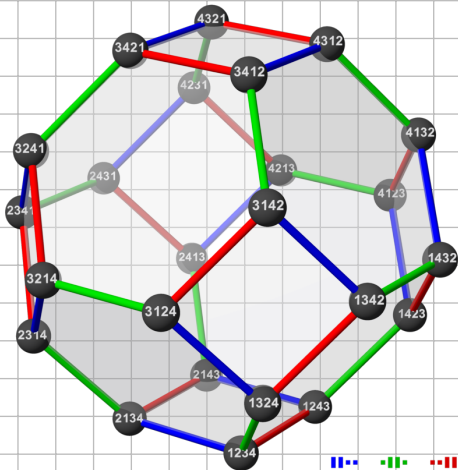
$a = (12)$
 $b = (23)$

$NC_3 = [1, (123)]$ inside \mathcal{O}_3
with generators



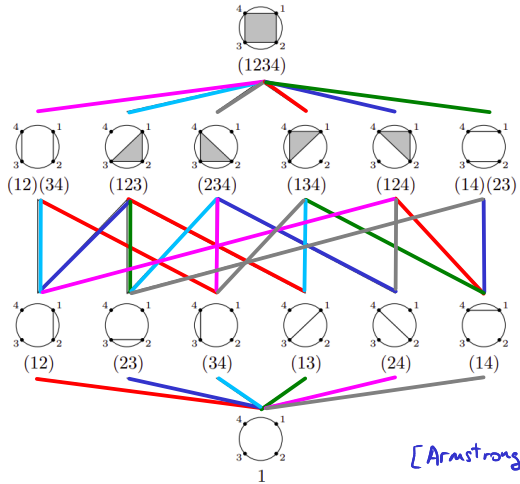
$a = (12)$
 $b = (23)$
 $c = (13)$

$(\mathcal{O}_4, \text{weak order})$



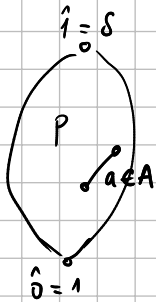
[Wikipedia]

$NC_4 = [1, (1234)]$ inside \mathcal{O}_4



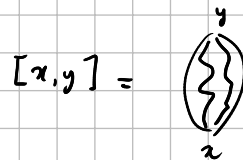
[Armstrong]

Let P be a finite bounded edge-labeled poset:



$A =$ set of labels for edges of P (actually appearing at least once)

$M = \langle A \mid \lambda(\gamma_1) = \lambda(\gamma_2) \text{ for any two maximal chains } \gamma_1, \gamma_2 \text{ in some } [x, y] \rangle_+$

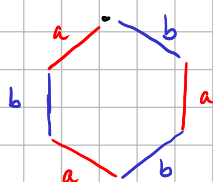


$\gamma = \{ x_0 < x_1 < \dots < x_k \}$
 $\lambda(\gamma) = a_1 a_2 \dots a_k \in A^*$

Monoid presentation

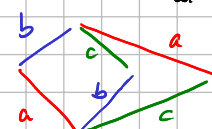
$G =$ group w/ same presentation

Example $(\mathcal{O}_3, \text{weak order})$



$a = (12)$
 $b = (23)$

$NC_3 = [1, (123)]$ inside \mathcal{O}_3
with generators



$a = (12)$
 $b = (23)$
 $c = (13)$

$$G = \langle a, b \mid aba = bab \rangle$$

braid group on 3 strands

$$a = \begin{array}{|c|} \hline \diagup \quad \diagdown \\ \hline \end{array} \quad b = \begin{array}{|c|} \hline \diagdown \quad \diagup \\ \hline \end{array}$$

$$G = \langle a, b, c \mid ab = bc = ca \rangle$$

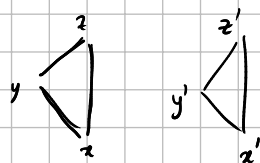
↓
the same group!

Theorem Let P be a combinatorial Garside structure. Then

- $P \hookrightarrow M \hookrightarrow G$
- There is a normal form $\delta^{m_1} x_1 \dots x_m$ with $x_i \in P$
- $K = \Delta(P) / \sim \cong K(G, 1)$
- Homological dim $(G) \leq l(P) \Rightarrow G$ torsion-free.

Given $x \leq y$ in P , let $\underline{d}(x, y) = \{ \lambda(\gamma) \mid \gamma \text{ maximal chain in } [x, y] \}$
"language of $[x, y]$ "

(C1) Group-like • Given $x \leq y \leq z$ and $x' \leq y' \leq z'$, if two pairs of languages are equal, then also the third pair of languages are equal



$$\underline{d}(x, y) = \underline{d}(x', y') \text{ and } \underline{d}(y, z) = \underline{d}(y', z') \\ \Rightarrow \underline{d}(x, z) = \underline{d}(x', z')$$

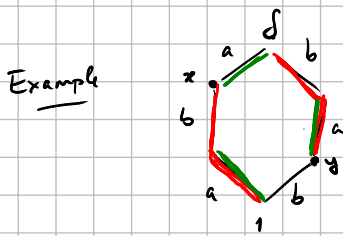
• $\forall x \in P, a \in A$, there is at most one edge $x < y$ labeled a
" " " $z < x$ " "



Sufficient condition: $P = [1, g]$ in the right Cayley graph of some group.

Remark: (C1) \Rightarrow if I know $\lambda(\gamma)$ for one chain of $[x, y]$, then $\underline{d}(x, y)$ is uniquely defined
(induction on $l(\gamma)$)
Exercise

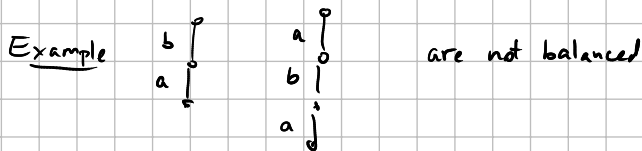
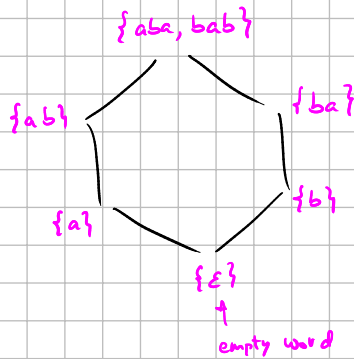
② Balanced $\{ \mathcal{L}(1, x) \mid x \in P \} = \{ \mathcal{L}(y, \delta) \mid y \in P \}$



- $\mathcal{L}(1, x) = \{ \underline{ab} \} = \mathcal{L}(y, \delta)$
- $\mathcal{L}(1, \delta) = \{ aba, bab \}$

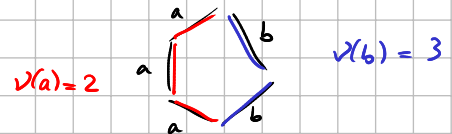
$= \{ \mathcal{L}(x, y) \mid x \leq y \text{ in } P \} =: L(P)$
 using ①
 (Exercise)

\Rightarrow We can identify $P \leftrightarrow L(P)$



Consequence of ② : There is a bijection $A \leftrightarrow$ atoms of P

③ Weakly graded There is a "norm" $\nu : A \rightarrow \mathbb{N}_+$ such that any two maximal chains in some $[x, y]$ have the same total norm



\Rightarrow Every language in $L(P)$ has a well-defined norm and also every element of \mathcal{M}

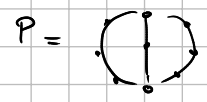
④ Lattice

Def If P satisfies ① - ④ then it is a combinatorial Garside structure, M is a Garside monoid, G is a Garside group.

Examples

- $P = B_n$ $M = \mathbb{N}^n$ $G = \mathbb{Z}^n$
- $P = (\mathcal{O}_n, \text{weak order})$ $M = \text{braid monoid}$ $G = \text{braid group on } n \text{ strands}$
- $P = NC_n$ $M = \text{dual braid monoid}$ $G \cong \uparrow$
- $G = \langle a_1, \dots, a_n \mid a_1^{m_1} = a_2^{m_2} = \dots = a_n^{m_n} \rangle$

Exercise: check (C1)-(C4)



• $\langle a_1, \dots, a_n \mid a_1 a_2 \dots a_k = a_2 a_3 \dots a_{k+1} = \dots = a_n a_1 a_2 \dots a_{n-1} \rangle$

Identity

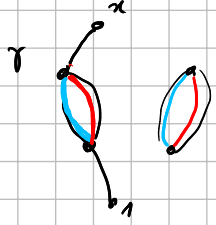
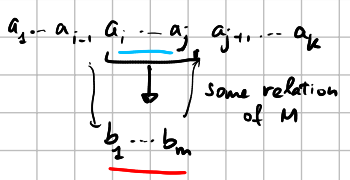
$P \cong L(P)$ so that $P \cong L(P) \rightarrow M \rightarrow G$
 $x \mapsto L(1, x)$

If all are injections, then P is an interval in the right Cayley graph of G .

Garside monoids

Lemma $P \rightarrow M$ is injective

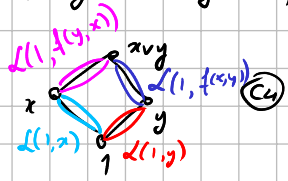
Proof Let $a_1 \dots a_k \in L(1, x)$ be a word that labels a chain γ in $[1, x]$ ($x \in P$).



By (C1), $a_1 \dots a_j$ determines its language

$\Rightarrow a_1 \dots a_{j-1} b_1 \dots b_m a_{j+1} \dots a_k$ labels a chain from 1 to $x \rightarrow$ gives the same element x at P \square

To define M : we can use $\bar{P} = P \setminus \{1\}$ as generators; given $x, y \in \bar{P}$,



$f: P \times P \rightarrow 1$
 $f(x, y) \in P$

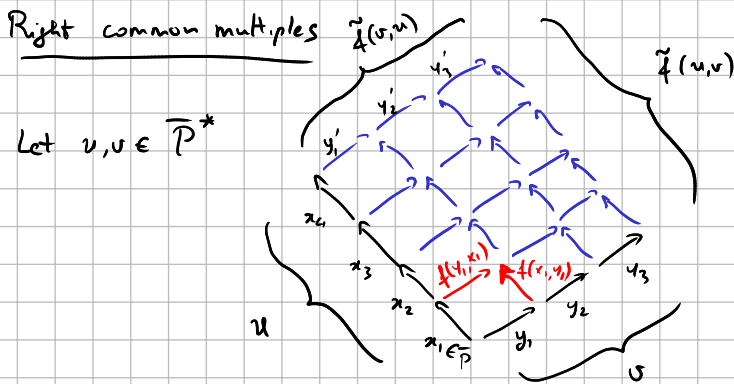
$L(x, xvy) = L(1, f(y, x))$
 $L(y, xvy) = L(1, f(x, y))$

$M = \langle \bar{P} \mid \underline{x} \underline{f(y, x)} = \underline{y} \underline{f(x, y)} \quad \forall x, y \in \bar{P} \rangle$

And these relations are enough because xvy is the join (Exercise)
 if $f(x, y) = 1$, then we put the empty word ϵ instead of $f(x, y)$



\Rightarrow Elements of M are represented by words in \bar{P} and two words $u, v \in \bar{P}^*$ are equivalent ($u \equiv v$) if they define the same element of M .



This defines $\tilde{f}: \bar{P}^* \times \bar{P}^* \rightarrow \bar{P}^*$
 extending f

By construction, $u \tilde{f}(v,u) \equiv v \tilde{f}(u,v)$

Lemma 1 Given $u, v, u', v' \in \bar{P}^*$, the following are equivalent:

- (i) $uu' \equiv vv'$
- (ii) $u' \equiv \tilde{f}(v,u)w$ and $v' \equiv \tilde{f}(u,v)w$ for some $w \in \bar{P}^*$.

Proof (ii) \Rightarrow (i) obvious
 (i) \Rightarrow (ii) a big induction [Dehornoy, Lemma 4]

Lemma 2 (a) The class of u in M left-divides the class of v iff $\tilde{f}(u,v) = \epsilon$. empty word

(b) $u \equiv v$ iff $\tilde{f}(u,v) = \tilde{f}(v,u) = \epsilon$ \rightsquigarrow we can check equality (\equiv) in M !

In particular, left divisibility in M is a partial order relation (antisymmetric)

Proof (a) $u \tilde{f}(v,u) \equiv v \tilde{f}(u,v)$ by the property of \tilde{f}

- If $\tilde{f}(v,u) = \epsilon$, then u left-divides v
- If u left-divides v , then $v\epsilon \equiv uu' \xrightarrow{\text{(Lemma 1)}} \epsilon \equiv \tilde{f}(u,v) \cdot w \xrightarrow{\text{(C3)}} \tilde{f}(u,v) = \epsilon$

(b). " \Rightarrow " follows from (a)

• " \Leftarrow " if $\tilde{f}(u,v) = \tilde{f}(v,u) = \epsilon$, then $u = u \frac{\tilde{f}(v,u)}{\epsilon} \equiv v \frac{\tilde{f}(u,v)}{\epsilon} = v$ □

Corollary 1 $v \tilde{f}(u,v) \equiv u \tilde{f}(v,u)$ is the right lcm of the classes of u and v in M .

\Rightarrow left gcd exist ($[1, uvv]$ is a finite join-semilattice (\leq is left divisibility) \Rightarrow lattice)

Corollary 2 M is left-cancellative: $uv \equiv u'v'$ and $u \equiv u' \Rightarrow v \equiv v'$

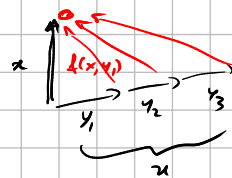
Proof $uv \equiv u'v' \xrightarrow{\text{Lemma 1}} \begin{cases} v \equiv \tilde{f}(u',u)w = w \\ v' \equiv \tilde{f}(u,u')w = w \end{cases} \Rightarrow v \equiv v'$ □
 Lemma 2

Corollary 3 (a) If the class of u left-divides some element $z \in P$, then it is in P

(b) $P \subseteq M$ is the set of left divisors of f .

Proof (a) We have $\tilde{f}(u, z) = \varepsilon$ by lemma 2.

By induction on $|u|$, this is a diagram in P
 $\Rightarrow u \in P$.



(b) It is enough to show that every $z \in P$ left-divides f in M (then we conclude with part (a))

This is obvious: $\tilde{f}(z, f) = f(z, f) = \varepsilon \xrightarrow{\text{lemma 2}} z \text{ left-divides } f \text{ in } M. \quad \square$

